

On the Complexity of Computing Real Radicals of Polynomial Systems

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Motivation

Polynomial system solving over the reals: $\mathbf{f} = (f_1, \dots, f_s) \in \mathbb{Q}[X_1, \dots, X_n]$

$$\mathbf{V}_{\mathbb{R}}(\mathbf{f}) = \{x \in \mathbb{R}^n \mid f_1(x) = 0, \dots, f_s(x) = 0\}$$

- ▶ Numeric computation \rightarrow reliability issues, especially in the singular case.
- ▶ Algebraic computation $\rightarrow (V = \mathbf{V}_{\mathbb{C}}(\mathbf{f}) \leftrightarrow \sqrt{\langle \mathbf{f} \rangle})$

What if $V \cap \mathbb{R}^n \subset \text{Sing}(V)$?

$$x_1^2 + x_2^2 = 0$$



$\mathbf{V}_{\mathbb{C}}(\mathbf{f})$: lines

$$\begin{cases} x_1 + ix_2 = 0 \\ x_1 - ix_2 = 0 \end{cases}$$

$\mathbf{V}_{\mathbb{R}}(\mathbf{f})$: point

$$x_1 = x_2 = 0$$

- Dimensions are different.
- $(0, 0) \in V$, singular
- $(0, 0) \in V \cap \mathbb{R}^n$, smooth!

Example 1 [Everett, Lazard, Lazard, Safey El Din, 2007]

$$\begin{aligned} \text{Vor1} = & (\alpha^2 + \beta^2 + 1)a^2\lambda^4 - 2a(2a\beta^2 + ay\beta + a\alpha x - \beta\alpha + 2a + 2a\alpha^2 - \beta\alpha a^2)\lambda^3 \\ & + (\beta^2 + 6a^2\beta^2 - 2\beta xa^3 - 6\beta\alpha a^3 + 6y\beta a^2 - 6a\beta\alpha - 2a\beta x + 6\alpha xa^2 + y^2 a^2 \\ & - 2a\alpha y + x^2 a^2 - 2y\alpha a^3 + 6a^2\alpha^2 + a^4\alpha^2 + 4a^2)\lambda^2 \\ & - 2(xa - ya^2 - 2\beta a^2 - \beta + 2a\alpha + \alpha a^3)(xa - y - \beta + a\alpha)\lambda \\ & + (1 + a^2)(xa - y - \beta + a\alpha)^2. \end{aligned}$$

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► Real zeros of Vor1 are union of:

$$\begin{cases} a\alpha - ax + \beta - y = 0 \\ \lambda + 1 = 0 \end{cases} \quad \begin{cases} a\alpha + ax - \beta - y = 0 \\ \lambda = 0 \end{cases} \quad \begin{cases} 2\beta\lambda + \beta + y = 0 \\ a = 0 \end{cases}$$

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- Only one connected component, which is not easy to be seen from Vor1.

Problem

- ▶ $\mathbf{f} = (f_1, \dots, f_s) \in \mathbb{Q}[X_1, \dots, X_n]$.
- ▶ $\sqrt[\text{re}]{\langle \mathbf{f} \rangle}$: the vanishing ideal of $\mathbf{V}_{\mathbb{R}}(\mathbf{f})$.
- ▶ An ideal I is called real if $I = \sqrt[\text{re}]{I}$.
- ▶ $D = \max\{\deg f_i, \dots, \deg f_s\}$.

Input: $\mathbf{f} = (f_1, \dots, f_s)$

Output: irreducible components of $\sqrt[\text{re}]{\langle \mathbf{f} \rangle}$:

- ▶ generators, or
- ▶ rational parametrizations.

Example 1 (Continued)

$$\begin{aligned}\text{Vor1} = & (\alpha^2 + \beta^2 + 1)a^2\lambda^4 - 2a(2a\beta^2 + ay\beta + a\alpha x - \beta\alpha + 2a + 2a\alpha^2 - \beta\alpha a^2)\lambda^3 \\ & + (\beta^2 + 6a^2\beta^2 - 2\beta xa^3 - 6\beta\alpha a^3 + 6y\beta a^2 - 6a\beta\alpha - 2a\beta x + 6\alpha xa^2 + y^2 a^2 \\ & - 2a\alpha y + x^2 a^2 - 2y\alpha a^3 + 6a^2\alpha^2 + a^4\alpha^2 + 4a^2)\lambda^2 \\ & - 2(xa - ya^2 - 2\beta a^2 - \beta + 2a\alpha + \alpha a^3)(xa - y - \beta + a\alpha)\lambda \\ & + (1 + a^2)(xa - y - \beta + a\alpha)^2.\end{aligned}$$

Irreducible components of $\sqrt[r]{\langle \text{Vor1} \rangle}$:

$$P_1 = \langle a\alpha - ax + \beta - y, \lambda + 1 \rangle$$

$$P_2 = \langle a\alpha + ax - \beta - y, \lambda \rangle$$

$$P_3 = \langle 2\beta\lambda + \beta + y, a \rangle$$

Timing: 9 sec.

State of the art

Exact computation:

- ▶ Becker, Neuhaus'1993, Neuhaus'1998, Spang'2007
Using **Gröbner bases** to compute real radicals for **arbitrary** polynomial ideals.
The complexity is $D^{2^{O(n^2)}}$.

Numerical approximations:

- ▶ Lasserre, Laurent, Rostalski'2008; Lasserre, Laurent, Mourrain, Rostalski, Trébuchet'2013
Using **SDP** relaxations to compute **zero-dimensional** real radical ideals.
- ▶ Ma, Wang, Zhi'2014
A **certificate** for computing real radicals using SDP relaxations.
- ▶ Brake, Hauenstein, Liddell'2016
A method based **SDP** programming for deciding if an ideal is real.

Main Results

$\mathbf{f} = (f_1, \dots, f_s) \subset \mathbb{Q}[X_1, \dots, X_n]$, $r = \dim \langle \mathbf{f} \rangle$, $D = \max\{\deg f_i\}$.

State of the art: $D^{2^{O(n^2)}}$

Smooth case.

A probabilistic algorithm computes **generators** of irreducible components of ${}^{re}\sqrt{\langle \mathbf{f} \rangle}$ using $(snD^n)^{O(1)}$ operations in \mathbb{Q} .

General case.

A probabilistic algorithm computes **rational parametrizations** of irreducible components of ${}^{re}\sqrt{\langle \mathbf{f} \rangle}$ using $s^{O(1)}(nD)^{O(nr2^r)}$ arithmetic operations in \mathbb{Q} .

Main idea

Simple point criterion [Bochnak, Coste, Roy, 1998]

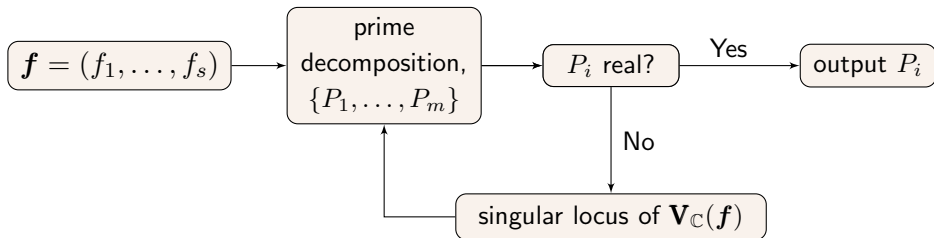
Prime $I = \langle f_1, \dots, f_s \rangle$ real $\iff \exists x \in \mathbf{V}_{\mathbb{R}}(I)$ s.t. $\text{rank} \left(\frac{\partial f_i}{\partial X_j}(x) \right) = n - r$,
where $r = \dim I$.

Main idea

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Main idea:



Singular point

Singular point [Cox, Little, O'Shea, 1992]

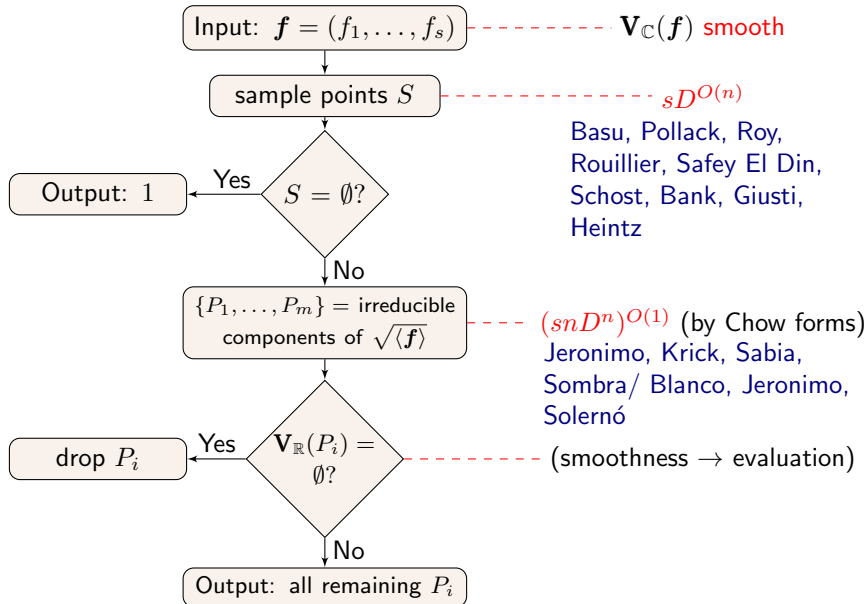
$V \subset \mathbb{C}^n$, $p \in V$, $\mathbf{I}(V) = \langle f_1, \dots, f_s \rangle$. The *tangent space* of V at p is

$$T_p(V) = \bigcap_{j=1}^s \left\{ x \in \mathbb{C}^n \mid \sum_{i=1}^n \frac{\partial f_j}{\partial X_i}(p) x_i = 0 \right\}.$$

$\dim_p V = \max\{\dim V_i \mid p \in V_i \text{ irreducible component of } V\}$.

- ▶ **Smooth Point:** $\dim T_p(V) = \dim_p V$.
- ▶ **Singular Point:** $\dim T_p(V) \neq \dim_p V$.
- ▶ **Singular locus:** $\text{Sing}(V) = \{p \in V \mid p \text{ is a singular point of } V\}$.
- ▶ V is **smooth** if $\text{Sing}(V) = \emptyset$.

Smooth case



General case

Drop the **smoothness assumption** on $V = \mathbf{V}_{\mathbb{C}}(\mathbf{f})$.

Difficulties: it may happen

- ▶ $V \cap \mathbb{R}^n \subset \text{Sing}(V)$;
- ▶ ... or even worse, in the **singular locus** of $\text{Sing}(V)$;

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Standard idea: lazy representations for equidimensional components of V .

- ▶ equations and **inequations**.
- ▶ Triangular set decompositions (Wu, Lazard, etc.) or **rational parametrizations** (Giusti, Heintz, Morais, Pardo, etc.)

Rational parametrization

An r -equidimensional variety $V \subset \mathbb{C}^n$ is the Zariski closure of the projection of the following set to $X = (X_1, \dots, X_n)$:

$$w(T) = 0, \quad X_i \frac{\partial w(T)}{\partial T_{r+1}} = v_i(T), \quad \frac{\partial w(T)}{\partial T_{r+1}} \neq 0$$

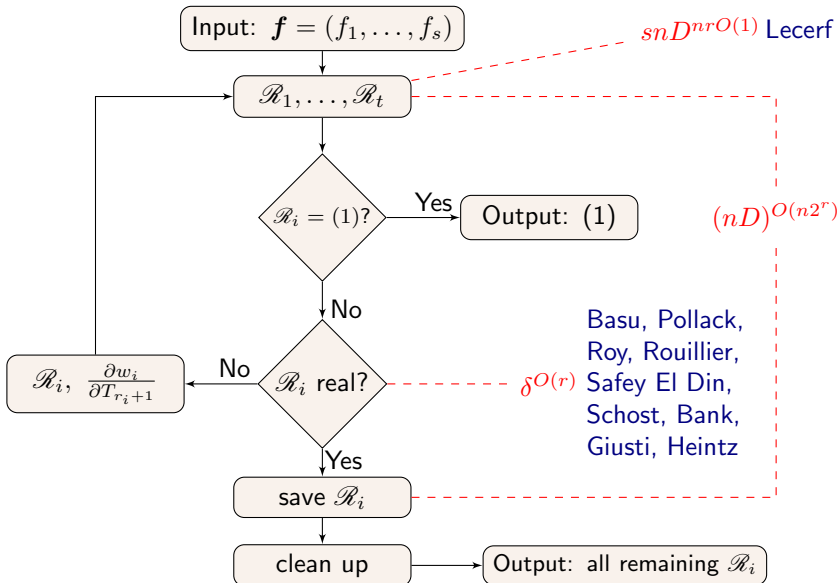
where $T = (T_1, \dots, T_{r+1})$, $i = 1, \dots, n$.

A rational parametrization of V :

- ▶ $\ell = (\lambda_1, \dots, \lambda_{r+1})$, generic linear combinations of X_1, \dots, X_n
- ▶ polynomials $w, v_1, \dots, v_n \in \mathbb{Q}[T]$.

Denote $\mathcal{R} = ((w, v_1, \dots, v_n), \ell)$.

General case



Implementation

Combing:

- ▶ SINGULAR: operating ideals [Greuel, Pfister].
- ▶ Maple: computing sample points RAGlib [Safey El Din] (uses FGb [Faugère] for computing Gröbner bases).

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Examples beyond the reach of the SINGULAR library realrad [Spang].

- ▶ (Homotopy-1) 7 variables, degree 7. [Chen, Davenport, May, Moreno Maza, Xia, Xiao]

$$f_1 = x^3y^2 + c_1x^3y + y^2 + c_2x + c_3, \quad f_2 = c_4x^4y^2 - x^2y + y + c_5, \quad f_3 = c_4 - 1.$$

Timing: 1 sec.

- ▶ (Essential variety) 9 variables, degree 3 [Fløystad, Kileel, Ottaviani]

$$\mathcal{E} = \{M \in \mathbb{R}^{3 \times 3} \mid \det(M) = 0, \quad 2(MM^T)M - \text{tr}(MM^T)M = 0\}$$

Timing: 800 sec.

Thank you!