

# Minkowski Decomposition and Geometric Predicates in Sparse Implicitization

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# Outline

- ① Implicitization by interpolation.
- ② Minkowski decomposition of the predicted polytope.
- ③ Geometric predicates as matrix operations.

## Implicitization by interpolation

# Implicitization

Given parameterization

$$x_0 = \alpha_0(t), \dots, x_n = \alpha_n(t), \quad t := (t_1, \dots, t_n),$$

compute the smallest algebraic variety containing the closure of the image of

$$\alpha : \mathbb{R}^n \rightarrow \mathbb{R}^{n+1} : t \mapsto \alpha(t), \quad \alpha := (\alpha_0, \dots, \alpha_n).$$

This is contained in the variety defined by the ideal

$$\langle p(x_0, \dots, x_n) \mid p(\alpha_0(t), \dots, \alpha_n(t)) = 0, \forall t \rangle.$$

When this is a *principal* ideal we wish to compute its defining polynomial  $p(x)$ , **given its Newton polytope**  $N(p(x))$  (*implicit polytope*), or, a superset of the exponents of its monomials with nonzero coefficient (*implicit support*).

# Approach

## Support prediction followed by interpolation.

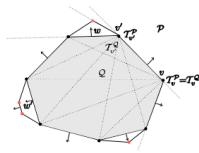
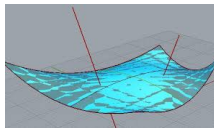
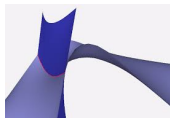
- Superset  $S$  of the implicit support: support of specialized sparse resultant.
- Construct  $\mu \times |S|$  ( $\mu \geq |S|$ ) matrix  $M$ : columns indexed by monomials with exponents in  $S$ , rows indexed by values of  $t$  at which the monomials are evaluated.
- The (ideally) unique kernel vector of  $M$  contains the coefficients of these monomials in the implicit equation.

## Example (Folium of Descartes)

$$x_0 = \frac{3t^2}{t^3 + 1}, \quad x_1 = \frac{3t}{t^3 + 1}, \quad S = \{(0, 3), (3, 0), (1, 1)\} \Rightarrow$$

$$M = \begin{bmatrix} x_1^3(\tau_1) & x_0^3(\tau_1) & x_0(\tau_1)x_1(\tau_1) \\ x_1^3(\tau_2) & x_0^3(\tau_2) & x_0(\tau_2)x_1(\tau_2) \\ x_1^3(\tau_3) & x_0^3(\tau_3) & x_0(\tau_3)x_1(\tau_3) \\ x_1^3(\tau_4) & x_0^3(\tau_4) & x_0(\tau_4)x_1(\tau_4) \end{bmatrix} \Rightarrow p(x_0, x_1) = x_0^3 - 3x_0x_1 + x_1^3.$$

# Previous work



16  
variables

- Integration of matrix  $SS^T$  over each parameter  $t_1, \dots, t_n$ . Successively larger supports to capture sparseness [Corless-Giesbrecht-Kotsireas-Watt '00].
- Successively larger supports also used in [Dokken-Thomassen '03] in the setting of approximate implicitization.
- Tropical geometry [Sturmfels-Tevelev-Yu '07] leads to algorithms for the polytope of (specialized) resultants [Jensen-Yu '12].
- Our method developed in [Emiris-Kalinka-Konaxis-Luu Ba '13a, '13b].

# Implicitization reduced to elimination

Setup (for omitted details see paper)

Given  $x_i = \frac{\alpha_i(t)}{\beta_i(t)}$ ,  $i = 0, \dots, n$ , define polynomials in  $(\mathbb{R}[x_0, \dots, x_n])[t]$ :

$$f_i := x_i \beta_i(t) - \alpha_i(t), \quad i = 0, \dots, n,$$

with supports  $A_i \subset \mathbb{Z}^n$ . If the parameterization is generically 1-1, then

$$p(x_0, \dots, x_n) = \text{Res}(f_0, \dots, f_n), \quad \text{provided that } \text{Res}(f_0, \dots, f_n) \neq 0.$$

## Projection

Let  $F := \{F_0, \dots, F_n\} \in \mathbb{R}(c_{ij})[t]$  be generic polynomials wrt  $A_i$  and let  $\phi$  be the specialization of their symbolic coefficients  $c_{ij}$  to those of the  $f'_i$ s:

$$\phi : c_{ij} \mapsto a_{ij} + b_{ij}x_i.$$

The  $\phi$ -projection of the Newton polytope  $N(\text{Res}(F))$  of the sparse resultant of polynomials in  $F$ , contains (a translate of) the Newton polytope of the implicit polynomial, or *implicit polytope*.

## Example

Given parameterization (Buchberger'88):

$$x_0 = t_1 t_2, \quad x_1 = t_1 t_2^2, \quad x_2 = t_1^2,$$

we define:

$$\begin{aligned} f_0 &:= x_0 - t_1 t_2, & f_1 &:= x_1 - t_1 t_2^2, & f_2 &:= x_2 - t_1^2, \\ F_0 &:= c_{00} - c_{01} t_1 t_2, & F_1 &:= c_{10} - c_{11} t_1 t_2^2, & F_2 &:= c_{20} - c_{21} t_1^2 \end{aligned}$$

$$\text{and } \phi(c_{00}, c_{01}, c_{10}, c_{11}, c_{20}, c_{21}) = (x_0, -1, x_1, -1, x_2, -1).$$

$$\text{Res}(F) = -c_{00}^4 c_{11}^2 c_{21} + c_{01}^4 c_{10}^2 c_{20}, \quad N(\text{Res}(F)) = ((4, 0, 0, 2, 0, 1), (0, 4, 2, 0, 1, 0))$$

and projection by  $\phi$  yields implicit polytope  $((4, 0, 0), (0, 2, 1))$ .

We construct  $M$  by evaluating  $(t_1, t_2)$  at random  $\tau_1, \tau_2, \tau_3 \in \mathbb{C}^2$ :

$$M = \begin{bmatrix} x_0^4(\tau_1) & x_1^2(\tau_1)x_2(\tau_1) \\ x_0^4(\tau_2) & x_1^2(\tau_2)x_2(\tau_2) \\ x_0^4(\tau_3) & x_1^2(\tau_3)x_2(\tau_3) \end{bmatrix}.$$

Kernel vector  $(-1, 1)$  yields implicit equation:  $-x_0^4 + x_1^2 x_2$ .

# Implicit support

- Input:  $n + 1$  supports  $A_i \subset \mathbb{Z}^n$ , and projection  $\phi$ .  
Output (predicted implicit polytope): the  $\phi$ -projection of the Newton polytope of the sparse resultant of  $A_0, \dots, A_n$ .
- ResPol [[Emiris-Fisikopoulos-Konaxis-Peñaranda'12](#)]  
computes precise polytope of bicubic surface in 1sec, with 715 terms.



- The lattice points in the implicit polytope correspond to the exponent vectors of the implicit support.

# Geometry of predicted support & higher dimensional kernel

## Assumption

The interpolation matrix  $M$  is built with sufficiently generic evaluation points.

## Theorem (Emiris-Kalinka-Konaxis-Luu Ba '13)

For projection  $\phi : F_i \mapsto f_i$  s.t. not all leading coefficients of  $f_i$  vanish,

$$\phi(\text{Res}(F)) = c(x) \cdot p(x).$$

If  $P = N(p)$  is the implicit and  $Q = \phi(N(\text{Res}(F)))$  the predicted polytopes, then

$$Q \supseteq E + P, \quad \text{for some polytope } E. \quad (+ : \text{Minkowski addition}).$$

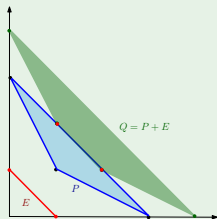
In particular, the kernel of  $M$  has dimension equal to the  $\#$  lattice points in  $E$ .

## Corollary

If  $v_1, \dots, v_k$  is a basis of the kernel of  $M$ , and  $g_1, \dots, g_k$  the corresponding polynomials, i.e.  $g_i = v_i^T S$ , then  $\text{gcd}(g_1, \dots, g_k) = p(x)$ ; can also use factoring

# Geometry of predicted support & higher dimensional kernel

**Example** (Folium of Descartes:  $x_0 = 3t^2/t^3 + 1$ ,  $x_1 = 3t/t^3 + 1$ )



Polynomials from kernel vectors

$$g_1 = x_0(x_0^3 - 3x_0x_1 + x_1^3)$$

$$g_2 = x_1(x_0^3 - 3x_0x_1 + x_1^3)$$

To extract the implicit polytope we employ Minkowski decomposition of  $Q$ .

## Minkowski decomposition of the predicted polytope

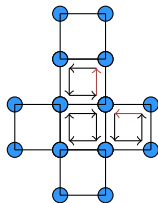
# Minkowski Decomposition of the Predicted polytope: Representation

## Input

A list of polygons: Facets of the 3-dimensional predicted polytope.

## Primitive edges

- $V = [v_1, v_2, \dots, v_m]$  ordered list of vertices.
- For every edge  $(v_{i_1}, v_{i_2})$ , define vector  $v_{i_2} - v_{i_1}$ .
- Let  $\ell_i$  be its integer length and  $e_i$  the corresponding primitive vector.



A cube flattened.  
Orientation shown in some of its facets. Common edges can have different orientation (red) in each facet.

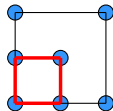
# The Integer Linear Program (ILP)

## Conditions for Decomposition

Given a facet  $F$ , it satisfies

$$\sum_{i \text{ s.t. } e_i \in F} \sigma_{i,F} \ell_i e_i = 0,$$

where  $\sigma_{i,F}$  is the sign of  $e_i$  and depends on the orientation of the edge  $\ell_i e_i$  in the facet  $F$ .



## ILP

$$a_i \in \mathbb{N}, \quad 0 \leq a_i \leq \ell_i,$$
$$\sum_{i \text{ s.t. } e_i \in F} \sigma_{i,F} a_i e_i = 0, \quad \text{for every facet } F.$$

Solutions that give summands homothetic to the input are excluded by adding two more appropriate equations.

A set of sub-edges  $a_i e_i$  (red) of the square forming a polygon.

# Recursion and Balance

## Balancing

We can choose the objective function of the ILP in order to decompose to polytopes with specific characteristics.

Due to the nature of the motivating problem, we prefer the decomposition to two summands of almost equal edge length sum, i.e., “balanced”.

This is achieved by setting as the objective function the sum of the edge lengths.

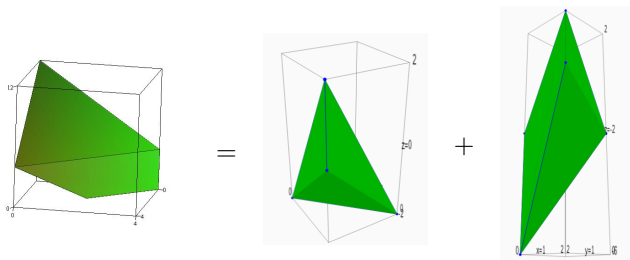
## Recursion

We recurse over the obtained two summand polytopes until we reach an indecomposable polytope.

## Example (Eight-surface)

$$x_0 = \frac{4s(-1 + t_2^2)(-1 + t_1^2)}{(1 + t_2^2)(1 + t_1^2)^2}, \quad x_1 = \frac{-8t_1 t_2(-1 + t_1^2)}{(1 + t_2^2)(1 + t_1^2)^2}, \quad x_2 = \frac{2t_1}{(1 + t_1^2)}$$

Predicted polytope leads to a  $67 \times 67$  interpolation matrix. Minkowski decomposition gives true implicit polytope (2nd summand) and a  $10 \times 10$  matrix.



## Geometric predicates as matrix operations

# Geometric predicates using interpolation matrices

## Setup

- $S$ : superset of implicit support,  
 $m(x)$ : vector of monomials (in the variables  $x_i$ ) with exponents in  $S$ .
- Fix *generic* distinct values  $\tau_k$ ,  $k = 1, \dots, |S| - 1$ .
- Construct  $(|S| - 1) \times |S|$  *numeric* matrix  $M'$  whose  $k$ th row,  $k = 1, \dots, |S| - 1$ , is vector  $m(x)$  evaluated at  $\tau_k$ .
- Construct

$$M(x) = \begin{bmatrix} M' \\ m(x) \end{bmatrix}.$$

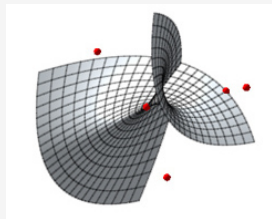
## Lemma

*Assuming  $M'$  is of full rank (equiv., the predicted polytope  $Q$  contains only one translate of  $P$ ), the determinant of matrix  $M(x)$  equals the implicit polynomial  $p(x)$  up to a constant.*

# Membership predicate

## Membership

Given  $x_i = f_i(t)/g_i(t)$ ,  $i = 0, \dots, n$ , and query point  $q \in \mathbb{R}^{n+1}$ , **decide if  $p(q) = 0$** , where  $p(x)$  is the implicit polynomial, by using the interpolation matrix.



## Lemma

Given  $M(x)$  and a query point  $q$  in  $(\mathbb{R}^*)^{n+1}$ , let  $M(q)$  denote matrix  $M(x)$  where its last row is evaluated by  $q$ . Then  $q$  lies on the hypersurface defined by  $p(x) = 0$  if and only if

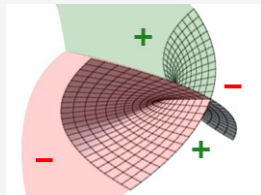
$$\text{corank}(M(q)) = \text{corank}(M').$$

*Numeric matrix  $M'$  need not be of full rank.*

# Sidedness predicate

Given a hypersurface in  $\mathbb{R}^{n+1}$  defined by  $p(x)$ , and point  $q \in \mathbb{R}^{n+1}$  such that  $p(q) \neq 0$ , we define

$$\text{side}(q) = \text{sign}(p(q)) \in \{-1, 1\}.$$



## Lemma

Assuming  $Q$  contains only one translate of  $P$ ,  $M(x)$  is the interpolation matrix and  $q \in (\mathbb{R}^*)^{n+1}$  such that  $p(q) \neq 0$ , then  $\det M(q) \neq 0$ .

## Lemma

Let  $q_1, q_2$  be two query points in  $(\mathbb{R}^*)^{n+1}$  not lying on the hypersurface  $p(x) = 0$ . Assuming  $Q$  contains only one translate of  $P$ , then

$$\text{side}(q_1) = \text{side}(q_2) \text{ iff } \text{sign}(\det M(q_1)) = \text{sign}(\det M(q_2)).$$

# Implementation

- Minkowski decomposition implemented in Sage.
- Sparse implicitization and geometric predicates implemented in Maple.
- Exact and approximate computations available.
- Option to use normal vectors to the curve/surface to build the interpolation matrix (Hermite interpolation).
- Code available at

<http://ergawiki.di.uoa.gr/index.php/Implicitization>

# Experimental comparison in Maple

Surface	Deg.	Msize	SI	GB	DR	Memb.	Side
Handkerchief surface	3	10	0.008	0.012	0.004	0.008	0.004
Moebius strip	3	398	16.784	0.268	0.204	15.012	83.0*
Bohemian dome	4	55	0.108	0.092	0.060	0.092	0.260
Eight surface	4	67	0.224	0.372	0.540	1.036	0.620
Swallowtail surface	5	25	0.032	0.032	0.008	0.048	0.044
Sine surface	6	125	0.556	0.160	0.048	0.532	7.916
Enneper's surface	9	103	0.268	0.204	0.032	0.476	3.468
Bicubic	18	715	16.804	>3000	8.028	68.176	1676.8*
HyperSurface	Deg.	Msize	SI	GB	DR	Memb.	Side
Bourgain hypersurface	3	6	0.004	0.0016	0.004	0.004	0.008
Hypercone	2	165	89.86	0.028	0.392	0.892	10.932*
Hypesurface	12	169	18.464	0.792	1.164	1.620	0.280

- Exact computations (Intel i5-2500, 3.30 GHz, 8GB Linux machine): runtimes (s) of: our method (SI), Gröbner bases (EliminationIdeal()) (GB), and Dixon resultant (Minimair) (DR).
- Last 2 columns show the runtime to decide Membership and Sidedness using previous Lemmas.
- Nullspace and determinants use Maple's native routines except \*: modular det computation.
- Experiments show our method is competitive and less dependent on the degree and the dimension of the (sparse) input. Matrices often ill-conditioned.



- Numerical computation: stability problems, non-rational coefficients, accuracy (Hausdorff distance).
- Ray shoot? More operations?
- Use Bernstein basis (no conversion to monomial basis).
- Exploit the generalized Vandermonde structure of  $M$ : implies  $O^*((\dim M)^2)$ . Relies on multivariate polynomial interpolation at arbitrary evaluation points.
- Higher codimension: space curves?

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Thank you for your attention!